# Hadamard-Diagonalizable Graphs with Perfect Quantum State Transfer 

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## Hadamard matrices

- A Hadamard Matrix $H$ is an $n \times n$ matrix whose entries are all 1 or -1 and satisfies $H H^{T}=n l$ (equivalently, its rows and/or columns are mutually orthogonal).
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& H_{2}=\left[\begin{array}{cc}
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\end{array}\right], \\
& H_{4}=H_{2} \otimes H_{2}=\left[\begin{array}{cc}
H_{2} & H_{2} \\
H_{2} & -H_{2}
\end{array}\right]=\left[\begin{array}{cccc}
1 & 1 & 1 & 1 \\
1 & -1 & 1 & -1 \\
1 & 1 & -1 & -1 \\
1 & -1 & -1 & 1
\end{array}\right]
\end{aligned}
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1 & 1 & 1 & 1 \\
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\end{array}\right], \\
H_{2^{n+1}} & =H_{2} \otimes H_{2^{n}}=\left[\begin{array}{cc}
H_{2^{n}} & H_{2^{n}} \\
H_{2^{n}} & -H_{2^{n}}
\end{array}\right] \text { for all } n \geq 1 .
\end{aligned}
$$

## Adjacency matrix of a (weighted) graph

- The adjacency matrix of a (weighted) graph is the $n \times n$ matrix $A=\left(a_{j, k}\right)$ defined by

$$
a_{j, k}= \begin{cases}w(j, k) & \text { if } \mathrm{j} \text { and } \mathrm{k} \text { are adjacent } \\ 0 & \text { otherwise }\end{cases}
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- For example, the graph on the left below has adjacency matrix on the right:



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- For example, the graph on the left below has adjacency matrix on the right:


$$
A=\left[\begin{array}{llllll}
0 & 5 & 3 & 0 & 0 & 0 \\
5 & 0 & 0 & 4 & 0 & 0 \\
3 & 0 & 0 & 1 & 2 & 0 \\
0 & 4 & 1 & 0 & 0 & 1 \\
0 & 0 & 2 & 0 & 0 & 4 \\
0 & 0 & 0 & 1 & 4 & 0
\end{array}\right]
$$

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$$
\begin{aligned}
D & =\left[\begin{array}{llllll}
8 & 0 & 0 & 0 & 0 & 0 \\
0 & 9 & 0 & 0 & 0 & 0 \\
0 & 0 & 6 & 0 & 0 & 0 \\
0 & 0 & 0 & 6 & 0 & 0 \\
0 & 0 & 0 & 0 & 6 & 0 \\
0 & 0 & 0 & 0 & 0 & 5
\end{array}\right], \\
L & =\left[\begin{array}{cccccc}
8 & -5 & -3 & 0 & 0 & 0 \\
-5 & 9 & 0 & -4 & 0 & 0 \\
-3 & 0 & 6 & -1 & -2 & 0 \\
0 & -4 & -1 & 6 & 0 & -1 \\
0 & 0 & -2 & 0 & 6 & -4 \\
0 & 0 & 0 & -1 & -4 & 5
\end{array}\right]
\end{aligned}
$$

## Laplacian matrix of a (weighted) graph

The Laplacian matrix $L$ of any graph has many nice properties:

- Symmetric.
- Positive semidefinite (since it is diagonally dominant).
- Row sums 0. Equivalently...
> ( $1,1, \ldots, 1$ ) is an eigenvector with eigenvalue 0 .
- The multiplicity of the eigenvalue 0 is the number of connected components.


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## Perfect State Transfer

A graph with Laplacian $L$ exhibits perfect state transfer (PST) at time $t$ between vertices $v_{j}$ and $v_{k}$ if the $(j, k)$-entry of $e^{i t L}$ has magnitude 1.

```
- Since \(L\) is symmetric, \(e^{i t L}\) is unitary, so none of its entries
    have magnitude larger than 1. Also, if PST occurs between
    vertices \(v_{j}\) and \(v_{k}\) then all other entries in the \(j\)-th row and
    \(k\)-th column of \(e^{i t L}\) are 0 .
- Motivation: This means that after time \(t\), the quantum state
    \(|j\rangle\) evolves into the state \(e^{i t L}|j\rangle=|k\rangle\).
- Morally, this means we have transferred the quantum state
    from vertex \(v_{j}\) to vertex \(v_{k}\) of the graph (perfectly-without
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A graph with Laplacian $L$ is Hadamard-diagonalizable if there exists a Hadamard matrix $H$ such that $H^{T} L H$ is diagonal.
> - It is more convenient to diagonalize by a scaled Hadamard $U=\frac{1}{\sqrt{n}} H$ so that $U^{T} L U$ contains the eigenvalues of $L$ along its diagonal.

- For example, the square graph is Hadamard-diagonalizable:



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L=\left[\begin{array}{cccc}
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0 & 4 & 0 & 0 \\
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\end{array}\right]\left(\frac{1}{2} H_{4}\right)^{T} .
\end{aligned}
$$

## Hadamard-Diagonalizable Graphs

WLOG, $H$ can be assumed to have every entry in its first row and column equal to 1 .

- Barik-Fallat-Kirkland (2011) says if a graph G has integer weights and is Hadamard-diagonalizable then:
- $G$ is regular (i.e., the sum of the edge weights adjacent to each vertex is constant).
- The eigenvalues of its Laplacian are even integers.
- Integer weights can be extended to rational weights via scaling (changes the time of PST).


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## Relationship with Cubelike Graphs

Lemma (Hadamard-Diagonalizability and Cubelike Graphs)
An unweighted graph is diagonalizable by the standard Hadamard if and only if it is cubelike.
> - Hadamard diagonalizability is thus more general than being cubelike.
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## A Spectral Characterisation

Given a spectral decomposition of a Hadamard-diagonalizable graph, the following theorem shows that it is easy to determine whether or not PST occurs at time $\pi / 2$.


> Corollary: If PST occurs, then half of the eigenvalues are
> $0(\bmod 4)$ and the other half are $2(\bmod 4)$.

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## Theorem (PST from Spectrum)

Let $G$ be an integer-weighted graph that is diagonalizable by a Hadamard matrix $H=\left(h_{i, j}\right)$. Denote the eigenvalues of its Laplacian by $\lambda_{1}, \cdots, \lambda_{n}$. Then $G$ has PST between vertices $v_{j}$ and $v_{k}$ at time $\pi / 2$ if and only if, for each $\ell=1, \cdots, n$,

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\lambda_{\ell} \equiv 1-h_{j, \ell} h_{k, \ell}(\bmod 4)
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0 & 0 & 0 & 0 \\
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$$

since

$$
H_{4}=\left[\begin{array}{cccc}
1 & 1 & 1 & 1 \\
1 & -1 & 1 & -1 \\
1 & 1 & -1 & -1 \\
1 & -1 & -1 & 1
\end{array}\right]
$$

## A Spectral Characterisation

This result can help us create graphs with PST.

- E.g., There is no unweighted Hadamard-diagonalizable graph on 12 vertices that has PST. But...
- The (essentially unique) $12 \times 12$ Hadamard is
$H_{12}=\left[\begin{array}{cccccccccccc}1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & -1 & 1 & -1 & 1 & 1 & 1 & -1 & -1 & -1 & 1 & -1 \\ 1 & -1 & -1 & 1 & -1 & 1 & 1 & 1 & -1 & -1 & -1 & 1 \\ 1 & 1 & -1 & -1 & 1 & -1 & 1 & 1 & 1 & -1 & -1 & -1 \\ 1 & -1 & 1 & -1 & -1 & 1 & -1 & 1 & 1 & 1 & -1 & -1 \\ 1 & -1 & -1 & 1 & -1 & -1 & 1 & -1 & 1 & 1 & 1 & -1 \\ 1 & -1 & -1 & -1 & 1 & -1 & -1 & 1 & -1 & 1 & 1 & 1 \\ 1 & 1 & -1 & -1 & -1 & 1 & -1 & -1 & 1 & -1 & 1 & 1 \\ 1 & 1 & 1 & -1 & -1 & -1 & 1 & -1 & -1 & 1 & -1 & 1 \\ 1 & 1 & 1 & 1 & -1 & -1 & -1 & 1 & -1 & -1 & 1 & -1 \\ 1 & -1 & 1 & 1 & 1 & -1 & -1 & -1 & 1 & -1 & -1 & 1 \\ 1 & 1 & -1 & 1 & 1 & 1 & -1 & -1 & -1 & 1 & -1 & -1\end{array}\right]$


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1 & 1 & 1 & 1 & -1 & -1 & -1 & 1 & -1 & -1 & 1 & -1 \\
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1 & 1 & -1 & 1 & 1 & 1 & -1 & -1 & -1 & 1 & -1 & -1
\end{array}\right]
$$

## A Spectral Characterisation

To make a graph with PST between vertices $v_{1}$ and $v_{2}$, we construct a set of integer eigenvalues that are 0 or $2(\bmod 4)$ according to the second row of $\mathrm{H}_{12}$ :


One possible choice:
$0,18,24,18,24,24,24,18,18,18,12,18$.
Then we set $\Lambda=\operatorname{diag}(0,18,24,18,24,24,24,18,18,18,12,18)$


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Then we set $\Lambda=\operatorname{diag}(0,18,24,18,24,24,24,18,18,18,12,18)$ and

$$
L=\frac{1}{12} H_{12}^{T} \wedge H_{12}
$$

## A Spectral Characterisation

It is straightforward to calculate

| $L$ | $=\frac{1}{12} H_{12}^{\top} \wedge H_{12}$ |
| ---: | :--- |
|  | $=\left[\begin{array}{cccccccccccc}18 & 0 & -1 & -1 & -1 & -3 & -3 & -3 & -1 & -3 & -1 & -1 \\ 0 & 18 & -1 & -1 & -1 & -3 & -3 & -3 & -1 & -3 & -1 & -1 \\ -1 & -1 & 18 & -2 & -2 & 0 & -2 & 0 & -2 & -2 & -4 & -2 \\ -1 & -1 & -2 & 18 & -4 & 0 & 0 & -2 & -2 & -2 & -2 & -2 \\ -1 & -1 & -2 & -4 & 18 & -2 & -2 & 0 & -2 & 0 & -2 & -2 \\ -3 & -3 & 0 & 0 & -2 & 18 & -2 & -2 & 0 & -2 & -2 & -2 \\ -3 & -3 & -2 & 0 & -2 & -2 & 18 & -2 & -2 & -2 & 0 & 0 \\ -3 & -3 & 0 & -2 & 0 & -2 & -2 & 18 & -2 & -2 & -2 & 0 \\ -1 & -1 & -2 & -2 & -2 & 0 & -2 & -2 & 18 & 0 & -2 & -4 \\ -3 & -3 & -2 & -2 & 0 & -2 & -2 & -2 & 0 & 18 & 0 & -2 \\ -1 & -1 & -4 & -2 & -2 & -2 & 0 & -2 & -2 & 0 & 18 & -2 \\ -1 & -1 & -2 & -2 & -2 & -2 & 0 & 0 & -4 & -2 & -2 & 18\end{array}\right]$ |

This is (necessarily, by construction) the Laplacian of a graph with PST.

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\begin{aligned}
L & =\frac{1}{12} H_{12}^{\top} \wedge H_{12} \\
& =\left[\begin{array}{cccccccccccc}
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0 & 18 & -1 & -1 & -1 & -3 & -3 & -3 & -1 & -3 & -1 & -1 \\
-1 & -1 & 18 & -2 & -2 & 0 & -2 & 0 & -2 & -2 & -4 & -2 \\
-1 & -1 & -2 & 18 & -4 & 0 & 0 & -2 & -2 & -2 & -2 & -2 \\
-1 & -1 & -2 & -4 & 18 & -2 & -2 & 0 & -2 & 0 & -2 & -2 \\
-3 & -3 & 0 & 0 & -2 & 18 & -2 & -2 & 0 & -2 & -2 & -2 \\
-3 & -3 & -2 & 0 & -2 & -2 & 18 & -2 & -2 & -2 & 0 & 0 \\
-3 & -3 & 0 & -2 & 0 & -2 & -2 & 18 & -2 & -2 & -2 & 0 \\
-1 & -1 & -2 & -2 & -2 & 0 & -2 & -2 & 18 & 0 & -2 & -4 \\
-3 & -3 & -2 & -2 & 0 & -2 & -2 & -2 & 0 & 18 & 0 & -2 \\
-1 & -1 & -4 & -2 & -2 & -2 & 0 & -2 & -2 & 0 & 18 & -2 \\
-1 & -1 & -2 & -2 & -2 & -2 & 0 & 0 & -4 & -2 & -2 & 18
\end{array}\right]
\end{aligned}
$$

This is (necessarily, by construction) the Laplacian of a graph with PST

## A Spectral Characterisation

It is straightforward to calculate

$$
\begin{aligned}
L & =\frac{1}{12} H_{12}^{T} \Lambda H_{12} \\
& =\left[\begin{array}{cccccccccccc}
18 & 0 & -1 & -1 & -1 & -3 & -3 & -3 & -1 & -3 & -1 & -1 \\
0 & 18 & -1 & -1 & -1 & -3 & -3 & -3 & -1 & -3 & -1 & -1 \\
-1 & -1 & 18 & -2 & -2 & 0 & -2 & 0 & -2 & -2 & -4 & -2 \\
-1 & -1 & -2 & 18 & -4 & 0 & 0 & -2 & -2 & -2 & -2 & -2 \\
-1 & -1 & -2 & -4 & 18 & -2 & -2 & 0 & -2 & 0 & -2 & -2 \\
-3 & -3 & 0 & 0 & -2 & 18 & -2 & -2 & 0 & -2 & -2 & -2 \\
-3 & -3 & -2 & 0 & -2 & -2 & 18 & -2 & -2 & -2 & 0 & 0 \\
-3 & -3 & 0 & -2 & 0 & -2 & -2 & 18 & -2 & -2 & -2 & 0 \\
-1 & -1 & -2 & -2 & -2 & 0 & -2 & -2 & 18 & 0 & -2 & -4 \\
-3 & -3 & -2 & -2 & 0 & -2 & -2 & -2 & 0 & 18 & 0 & -2 \\
-1 & -1 & -4 & -2 & -2 & -2 & 0 & -2 & -2 & 0 & 18 & -2 \\
-1 & -1 & -2 & -2 & -2 & -2 & 0 & 0 & -4 & -2 & -2 & 18
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## Merge of Graphs

Suppose that $G_{1}$ and $G_{2}$ are two weighted graphs that are both diagonalizable by a Hadamard matrix $H$, with Laplacians $L_{1}=D_{1}-A_{1}$ and $L_{2}=D_{2}-A_{2}$, respectively.

We define their merge with respect to the weights $w_{1}$ and $w_{2}$ to be the graph $G_{1}{ }_{w_{1}} \odot_{w_{2}} G_{2}$ with Laplacian


When $w_{1}=w_{2}=1$, we denote the merge simply by $G_{1} \odot G_{2}$.

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\left[\begin{array}{cc}
w_{1} L_{1}+w_{2} D_{2} & -w_{2} A_{2} \\
-w_{2} A_{2} & w_{1} L_{1}+w_{2} D_{2}
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## Example



## Using the Merge to Create Graphs with PST

Our main interest in the merge comes from the fact that we can use it to create larger graphs with PST from smaller ones.


1) $j, k \in\{1, \ldots, n\}$ and
 time $\pi / 2$, or
 and $v_{k}$ at time $\pi / 2$, or
1.c) $w_{1}$ and $w_{2}$ are odd, $d_{2}$ is even, and the weighted graph with Laplacian $L_{1}+L_{2}$ has PST between $v_{j}$ and $v_{k}$ at time $\pi / 2$;

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## Theorem (Merge for PST)

Suppose $G_{1}$ and $G_{2}$ are integer-weighted graphs with regularities $d_{1}$ and $d_{2}$, respectively, both of which are diagonalizable by the same Hadamard matrix $H$. Then $G_{1}{ }_{w_{1}} \odot_{w_{2}} G_{2}$ has PST between vertices $v_{j}$ and $v_{k}$ at time $\pi / 2$ if and only if one of the following 8 conditions holds:


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## Using the Merge to Create Graphs with PST

```
2) j,k\in{n+1,\ldots,2n} and
    2.a) }\mp@subsup{w}{1}{}\mathrm{ is odd, w}\mp@subsup{w}{2}{}\mathrm{ is even, and G}\mp@subsup{G}{1}{}\mathrm{ has PST between }\mp@subsup{v}{j-n}{}\mathrm{ and }\mp@subsup{v}{k-n}{
        at time }\pi/2\mathrm{ , or
    2.b) }\mp@subsup{w}{1}{}\mathrm{ and }\mp@subsup{d}{2}{}\mathrm{ are even, w}\mp@subsup{w}{2}{}\mathrm{ is odd, and }\mp@subsup{G}{2}{}\mathrm{ has PST between }\mp@subsup{v}{j-n}{
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    Laplacian L}\mp@subsup{L}{1}{}+\mp@subsup{L}{2}{}\mathrm{ has PST between }\mp@subsup{v}{j-n}{}\mathrm{ and }\mp@subsup{v}{k-n}{}\mathrm{ at time }\pi/2\mathrm{ ;
3) }j\in{1,\ldots,n},k\in{n+1,\ldots,2n} and
    3.a) w
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    3.b) w
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3.b) $w_{1}, w_{2}$, and $d_{2}$ are all odd, and the weighted graph with
Laplacian matrix $L_{1}+L_{2}$ has PST between $v_{j}$ and $v_{k-n}$ at time $\pi / 2$.

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- If both $w_{1}$ and $w_{2}$ are even, the merge $G_{1} w_{w_{1}} \odot_{w_{2}} G_{2}$ does not have PST at time $\pi / 2$.
- However, it will have PST at some other time via re-scaling (divide by the highest common power of 2 factor of $w_{1}, w_{2}$ ).
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## Using the Merge to Create Graphs with PST

## Corollary (PST of Weighted Hypercube)

Suppose $w_{1}, w_{2}, \ldots, w_{n}$ are nonzero integers, exactly $d$ of which are odd, and consider the weighted hypercube
$C_{n}:=\left(w_{1} K_{2}\right) \square\left(w_{2} K_{2}\right) \square \cdots \square\left(w_{n} K_{2}\right)$. For each vertex $v_{j}$ of $C_{n}$, there is a vertex $v_{k}$ at distance $d$ from $v_{j}$ such that there is PST from $v_{j}$ to $v_{k}$ at time $\pi / 2$.

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## Open Questions

- Are there Hadamard-diagonalizable graphs with PST of all sizes that are a multiple of 4?
- Is there a Hadamard-diagonalizable graph with PST associated with each Hadamard matrix?
- If $G$ is unweighted, Hadamard diagonalizable, $r$-regular, has $n$ vertices, and has PST at time $\pi / 2$. Is it the case that $n \leq 2^{r}$ ?
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